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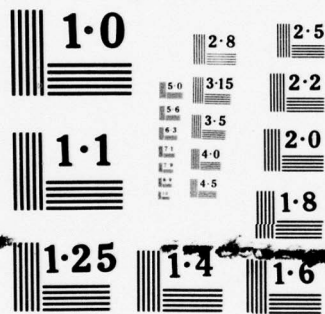
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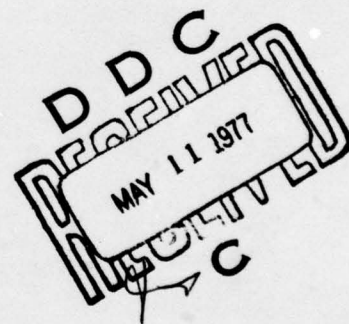
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# SOLVING THE REAL GENERALIZED SYMMETRIC EIGENPROBLEM BY SIMULTANEOUS ITERATION

ANALYSIS AND OPTIMIZATION BRANCH  
STRUCTURAL MECHANICS DIVISION

DECEMBER 1976

TECHNICAL REPORT AFFDL-TR-76-118  
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FOREWORD

This report describes work performed in the Aerospace Research Laboratories and in the Air Force Flight Dynamics Laboratory under the Defense Research Sciences Program, Project 7071, Research in Applied Mathematics, Task 01, Computational Aspects of Fluid and Structural Mechanics. The work described was carried out between October 1972 and July 1976. This is an interim report. Further reports in this series will be documented under Project 2304, Mathematical and Information Sciences, managed by the Air Force Office of Scientific Research, Bolling AFB, DC.

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## SECTION I

## INTRODUCTION

The structural analysis of mathematical models of complex aerospace and naval vehicles often involves the partial solution of the generalized algebraic eigenvalue problem. A pair of  $n$ -square matrices  $A$  and  $B$  are given and it is required to calculate one or more non-zero eigenvectors  $u$  and corresponding complex eigenvalues  $\lambda$  for which  $Au = \lambda Bu$ . In many familiar applications the matrices  $A$  and  $B$  are real and symmetric and  $B$  is positive definite so that all arithmetic is real. But  $n$  is very large. Often  $A$  and  $B$  have relatively few non-zero entries; i.e.,  $A$  and  $B$  are sparse. Nevertheless, the partial numerical solution of the large generalized symmetric eigenproblem is a challenge to the numerical analyst and the systems programmer alike. The terms "large" and "very large" are used here relative to the computer at hand and the limitations imposed by its operating system. One may distinguish three cases. If at most one copy of  $A$  and  $B$  can be stored in the central memory of the machine, we say that  $n$  is large. If  $n$  is so great that only a relatively few  $n$ -vectors may be simultaneously stored, then  $n$  is very large. If at most one  $n$ -vector can be accommodated, then  $n$  is simply gigantic.

This report describes the simultaneous iteration method for the partial solution of the generalized symmetric eigenproblem and gives a USA Standard FORTRAN implementation of the algorithm called SUBROUTINE SIMITZ. The use of SIMITZ is indicated if  $n$  is large or very large,



but it will not handle gigantic problems. However, it may be appropriate to problems of moderate size when only a few eigenvalues or eigenvectors are required.

Section II of this report provides a brief mathematical formulation of the simultaneous iteration method and describes the Rutishauser-Reinsch algorithm *ritzit* upon the program SIMITZ is based. Section III is devoted to a discussion of alternative methods for solution of the large eigenproblem and their comparison with simultaneous iteration. Detailed instructions for use of SIMITZ are given in Section IV and a complete listing of the FORTRAN program comprises Section V.



## SECTION II

## DESCRIPTION

The present program is an implementation of the simultaneous iteration algorithm [2] for calculating the eigenvalues largest in magnitude and corresponding eigenvectors of a real matrix symmetric relative to a prescribed inner product. Let  $ip(n, w, z)$  denote an inner product in the space of real column  $n$ -tuples and let the real  $n$ -square matrix  $C$  satisfy  $ip(n, Cw, z) = ip(n, w, Cz)$ . Then  $C$  is symmetric relative to  $ip$ , and if the  $n$ -square positive definite matrix  $B$  satisfies  $ip(n, w, z) = w^T Bz$  then  $C$  is B-symmetric. The equation  $BC = C^T B$  characterizes the B-symmetry of  $C$ . Given an optional set of  $p$  initial approximate eigenvectors of a real  $n$ -square B-symmetric matrix  $C$  corresponding to  $p$  eigenvalues of  $C$  largest in magnitude, the program calculates  $em$  eigenvalues and  $em$  corresponding eigenvectors,  $0 \leq em < p \leq n$ , to a precision dependent on the structure of  $C$  and on a prescribed tolerance  $eps$ . The matrix  $B$  is presented to the program as an independently prepared real function subprogram which calculates  $ip(n, w, z) = w^T Bz$  given column  $n$ -vectors  $w$  and  $z$ . The matrix  $C$  is presented as an independently prepared subroutine subprogram  $op(n, w, z)$  which when given an  $n$ -vector  $z$  computes its image  $w = Cz$ . The program is an outgrowth of a literal FORTRAN translation [6] of the ALGOL procedure *ritsit* [9] to which it is substantially equivalent when  $ip(n, w, z) = w^T z$ , the standard inner product. But depending on the choice of  $B$  and  $C$ , the present

program enables the direct treatment of a wide variety of symmetric eigenproblems.

Let  $A = A^T$  and  $B = B^T$  denote  $n$ -square real matrices and let  $\sigma$  be real. If  $B$  is positive definite then the matrix  $C = B^{-1}(A - \sigma B)$  is  $B$ -symmetric, and the program computes eigenvalues farthest from  $\sigma$  of the eigenproblem  $Au = \lambda Bu$  and corresponding eigenvectors. Implementation of  $op(n, w, z)$  here consists in providing for the appropriate solution for  $w$  of the linear system  $Bw = (A - \sigma B)z$ . Alternatively, selection of  $op$  to solve the system  $(A - \sigma B)w = Bz$  for  $w$  enables the calculation by simultaneous inverse iteration of the eigenvalues nearest to  $\sigma$  and their eigenvectors. Implications for large sparse systems for which the Cholesky factorization [7] of  $B$  is impractical are clear.

Let the eigenvalues  $d_1, \dots, d_p, d_{p+1}, \dots, d_n$  of  $C$  be arranged in order of descending absolute value and let  $E_p$  denote the direct sum of the distinct eigenspaces corresponding to  $d_1, \dots, d_p$ . Let  $X_0$  denote an  $n$ -by- $p$  matrix having a  $p$ -dimensional column space not orthogonal relative to  $ip$  to any eigenvector in  $E_p$ . Simultaneous iteration is based on the observation that if  $|d_p| > |d_{p+1}|$ , the columns of the matrix  $X_{k+m} = C^m X_k$  tend to a basis of  $E_p$  as  $k$  increases. But in practice all of the columns of  $X_{ks}$  tend toward the eigenspace  $E_1$  causing loss of information concerning the residual eigenvectors. To counter this tendency, set

$$X_{k+m} = C^m X_k B_{k+m}^{-1} \quad (1)$$

where the  $p$ -square upper triangular matrix  $B_{ks}$  is constructed together with  $X_{ks}$  by the Gram-Schmidt process to render the columns of  $X_{ks}$  orthonormal relative to  $ip$ . Now the  $i$ -th column vector of  $X_{ks}$  converges to the  $i$ -th eigenvector of  $C$  at a rate proportional to  $\max_{2 \leq i \leq p} (|d_i/d_{i-1}|, |d_{i+1}/d_i|)$ . Clearly this convergence will be delayed in the presence of eigenvalue clustering. But if  $|d_p| - |d_{p+1}|$  is not too small, the column space of  $X_{ks}$  will contain a good approximation to the  $i$ -th eigenvector even when  $ks$  is small.

In order to recover this approximation, a modified Rayleigh-Ritz process is employed. Let  $Q_{ks}$  denote an orthogonal matrix which diagonalizes the  $p$ -square symmetric matrix  $B_{ks} B_{ks}^T$ . Then the  $i$ -th column vector of

$$X_{k+1} = C X_k B_{k+1}^{-1} Q_{k+1} \quad (2)$$

converges to the  $i$ -th eigenvector of  $C$  at a rate proportional to  $|d_{p+1}/d_i|$  while the entries of the diagonal matrix computed with  $Q_{ks}$  and properly ordered offer close approximations to  $d_1^2, \dots, d_p^2$ . The true signed eigenvalues need only be computed at termination by diagonalizing the leading  $(p-1)$ -square principal submatrix of  $X_{ks}^T B C X_{ks}$ , the eigenproblem for  $C$  projected on  $E_{p-1}$  relative to  $ip$ .

The program determines a strategy for employing the devices (1) and (2) based on the distribution of the leading  $p$  eigenvalues of  $C$  upon which the convergence rate ultimately depends. The selection of values  $m$  in (1) is particularly important in this regard in that  $C^m X_k$  is replaced by the  $m$ -th Chebychev polynomial on an appropriate inter-



val  $[-e, e]$  evaluated by a special 3-term recurrence relation and permitting accelerated convergence when values of  $m$  are continually large. As a result the convergence quotient lies between  $|d_p/d_{em}|$  and  $\exp(-\text{arc cosh } |d_{em}/d_p|)$ . It is nearer to the first value if  $|d_1/d_{em}|$  is large and nearer to the second if the latter quotient is close to one.

As the iteration proceeds through a maximum of  $|km|$  iteration steps -  $km$  is a program parameter - acceptance tests for the eigenvalues and eigenvectors are conducted following each of the Rayleigh-Ritz steps (2). As soon as the relative increase of  $|d_{h+1}|$  is smaller than  $\text{eps}/10$ , then  $d_{h+1}$  is accepted and  $h$ , the number of previously accepted eigenvalues, is increased by one. Eigenvectors are accepted in groups of one or more corresponding to clusters of accepted eigenvalues nearly equal in magnitude. If  $g$  eigenvectors have already been accepted, let  $d_{g+1}, \dots, d_\ell$  denote such a cluster. For all  $j$ ,  $g+1 \leq j \leq \ell$ , denote by  $y_j$  the projection relative to  $ip$  of the image  $Cx_j$  of the  $j$ -th column  $x_j$  of  $X_{ks}$  on the linear closure of  $x_1, \dots, x_\ell$ . Set  $f_i = \max_j ||Cx_j - y_j|| / ||Cx_j||$  for  $i = g+1, \dots, \ell$  where the indicated norm is the Euclidean norm or 2-norm relative to  $ip$ . If  $|d_\ell|f_\ell / (|d_\ell| - e)$  is smaller than  $\text{eps}$  then all the  $x_j$ ,  $j = g+1, \dots, \ell$ , are accepted as eigenvectors and  $g$  is increased to  $\ell$ . The error quantities  $f_i$  are systematically discounted in accordance with the convergence properties of the algorithm to permit convergence in the presence of excessive round-off error or in case the parameter  $\text{eps}$  is prescribed unrealistically small. Having determined  $g$  eigen-

vectors, the iteration continues with  $p - g$  remaining columns of  $X_{ks}$  until either  $em$  eigenvectors have been calculated or  $|km|$  has been exceeded. The program may reduce  $em$  if it detects either no progress in convergence of eigenvectors corresponding to smaller eigenvalues or lack of stability in the behavior of larger eigenvalues.

The user may wish to supplement the forgoing outline of the operation of the program by consulting the description of the ALGOL procedure *ritzit* in [9] or [12] as well as a review of the mathematical foundations of simultaneous iteration in [5] and [8]. For this reason we describe the principal differences between *ritzit* and the present program. (a) The procedure *inprod* for calculating standard inner products was removed and the procedure *ip* was introduced where appropriate. (b) The procedure *jacobi* for calculating the solutions of the eigenproblem for the  $p$ -square and  $(p - 1)$ -square symmetric matrices was replaced by calls to the EISPACK [9] subroutines TRED2 and IMTQL2, primarily to save space. (c) The procedure *random* for calculating random column  $n$ -vectors of the matrix  $X_{ks}$  was replaced by in-line code which references a FORTRAN function RANF. RANF returns uniformly distributed random REAL values from the interval  $(0, 1)$ , one per function reference, given any one argument of any type. It is provided by the user. (d) The procedure *orthog* to perform Gram-Schmidt orthogonalization of the columns of  $X_{ks}$  was replaced by internally linked in-line code. In attempting to control potential underflow within *orthog* in a machine independent fashion,



*ritzit* calculates the machine precision  $mc$  but assumes in usage that out-of-range values underflow gracefully to zero, a machine dependent characteristic. The present program utilizes a single REAL machine dependent constant  $MT$ , the ratio of the smallest FORTRAN representable positive value to the machine precision, to test for this condition and upon its detection to take appropriate measures. (e) In its ALGOL implementation *ritzit* requires approximately  $(p + 3)n + 2p^2 + 5p$  storage locations in excess of those required by the program. Economies resulting in part from (b) above [3] have reduced this requirement to  $(p + 2)n + p^2 + 4p$  in the present program. All working storage is confined to a single array of  $2n + p^2 + 3p$  locations. (f) The value of  $km$  as an input parameter, set to  $|km|$  during program execution, is finally replaced by the value of  $ks$  as an output parameter, the number of iteration steps used in the calculation of  $em$  eigenvectors. (g) The present program retains unchanged the reference to a user supplied procedure *inf* as a window on program execution. However, the one variable involving  $eps$  is periodically redefined to enable effective control of  $eps$  from *inf* or from  $ip$  or  $op$  should this prove desirable.

The testing procedures developed for the present FORTRAN program parallel its evolution from a research tool, which conformed closely to its ALGOL parent, to present form. Early testing was concentrated on duplicating the tests furnished with *ritzit* and in eliminating errors in interpretation and translation of the ALGOL code. This was done for the most part on the CDC 6600 using FORTRAN Extended, Ver-

sion 3, under the SCOPE 3.3 operating system. Upon completion of this first phase, the resulting program was distributed as SUBROUTINE RITZIT with a locally developed library of FORTRAN linear algebra routines [6]. This same program served as a basis of *ritzit* translations for IBM 360/370 processors [1,3] whose preparation uncovered several bugs in the RITZIT code and suggested worthwhile modifications. A second phase of testing involved the development of a package of auxiliary FORTRAN programs for use with SUBROUTINE RITZIT to solve the eigenvalue problem  $Au = \lambda Bu$  through methods depending on Cholesky factorization where A and B may either be full matrices or sparse and banded. This phase was conducted on the CDC 6600 using FORTRAN Extended, Version 4, under SCOPE 3.4.

Systematic testing of the present program, SUBROUTINE SIMITZ, has been accomplished in part with the aid of driver program TESTB which generates a symmetric band matrix A and a lower triangular band matrix T of prescribed order n and bandwidths whose relevant entries are randomly generated integer values from a prescribed interval. The band matrix B is  $TT^T$ , and the program calculates the maximal eigenvalues of  $Au = \lambda Bu$ . For the sequence of values of p,  $p = 2, \dots, \min(\lfloor n/5 \rfloor, 10)$ , TESTB exercises SIMITZ for successive values of  $e_m$ ,  $e_m = 1, \dots, p - 1$ . For each value of i,  $i = 1, \dots, e_m$ , TESTB computes the residuals  $Ax_i - d_i Bx_i$  and their Euclidean norms relative to the standard inner product. Each norm is normalized by the difference  $|d_1| - e$ , and for each value of  $e_m$  the quantity  $\max_{1 \leq i \leq e_m} \|Ax_i - d_i Bx_i\| / (|d_1| - e)$ , the value k of i for which

the maximum occurs, and the corresponding geometric mean with unit weights are listed. Also listed are the relevant non-zero diagonals of A and T and the final eigenvalues computed for  $em = p - 1$ .

Figure 1 shows an output listing from the executable program TESTB on the CDC 6600 under the NOS/BE operating system and FORTRAN Extended, Version 4. Here A and B are of order 30 and each of bandwidth 7 having relevant entries between -99 and +99. Listed are the main diagonal and the three adjacent lower diagonals of T and A beginning with the entries in the first column. Here  $\epsilon = 10^{-10}$  and  $km = 100$ . Note how the relative nearness in magnitude of the first three eigenvalues inhibits the convergence of the second eigenvector when  $p = 3$  and  $em = 2$  resulting in acceptance of the first eigenvector only. The fourth eigenvalue, however, is in absolute value far enough away from this cluster to permit successful convergence when  $p = 4$  and  $em = 1, 2$ , and  $3$ . This phenomenon points to a procedure for pursuing a solution when  $p$  is initially chosen too small. SIMITZ may be reentered with X containing the approximate eigenvectors calculated for the smaller value of  $p$  as initial approximations for use with  $p$  increased in size. Significant processor time may often be saved in this fashion.

ADJACENT NON-ZERO LOWER DIAGONALS OF T, TT' = 3														
53	86	78	60	85	99	73	64	80	88	93	93	57	52	93
74	99	88	87	52	64	86	95	65	86	55	95	51	76	70
-26	32	49	65	-30	38	82	-9	-87	56	18	35	33	-46	-31
1	90	-44	75	-6	-49	99	90	-62	16	2	-40	-57	14	
-25	6	54	32	-73	-35	-1	-46	6	-13	-77	77	-51	-85	31
-24	88	-66	-18	-41	5	80	74	-80	-2	-6	78	-37		
-76	19	-26	-5	97	70	6	-95	45	-30	-12	-55	-92	9	39
37	-19	26	86	-48	-52	-31	49	-6	-85	5	21			
ADJACENT NON-ZERO LOWER DIAGONALS OF A = A'														
-3	50	-55	-73	-47	88	-33	-28	-77	-40	96	78	9	-54	-39
10	95	-53	-80		52	44		-50	-94	-22	29	-36	-4	36
0	65	59	-29	-84	-4	41	-91	72	-65	74	36	32	-86	-5
22	56	94	9	-47	84	-97	-1	57	-22	-1	-55	-15	51	
52	5	-41	-95	-47	-91	14	84	-83	83	43	-27	61	-79	14
21	85	-5	-61	-16	-15	90	-44	-22	-66	50	-34	-37		
-17	-84	78	-95	-89	-31	-43	52	-64	-70	71	95	-63	-12	87
91	-93	6	39	66	32	-21	92	-11		-98	-12			
										MEAN RESIDUAL				
										TIME (SECS)				
2	EM(IN)	EM(OUT)	KS	K	MAX RESIDUAL					MEAN RESIDUAL				
3	1	1	30	1	3.70E-11					3.70E-11				
3	1	1	26	1	2.81E-11					2.81E-11				
3	2	1	56	1	3.00E-11					3.00E-11				
4	1	1	15	1	2.17E-11					2.17E-11				
4	2	2	15	2	2.01E-11					1.91E-11				
4	3	3	24	3	2.53E-11					2.09E-11				
5	1	1	13	1	2.20E-11					2.20E-11				
5	2	2	13	1	2.03E-11					1.99E-11				
5	3	3	13	1	2.24E-11					1.50E-11				
5	4	4	33	1	2.37E-11					6.67E-12				
6	1	1	13	1	2.09E-11					2.09E-11				
6	2	2	13	2	2.16E-11					2.04E-11				
6	3	3	13	2	2.89E-11					2.03E-11				
6	4	4	21	4	2.02E-11					3.63E-11				
6	5	4	31	1	2.67E-11					1.09E-11				

## FINAL EIGENVALUES

2.393019487013E+00 -1.245936323012E+00 1.188248083996E+00 -2.273759446561E-01 -9.7.5362962256E-02

Figure 1. Output Listing

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## SECTION III

## SOME ALTERNATIVE METHODS

The generalized symmetric eigenvalue problem  $Au = \lambda Bu$  may in theory be reduced to a standard symmetric eigenvalue problem  $Pv = \lambda v$ . When  $B$  is positive definite a lower triangular matrix  $T$  may be calculated by Cholesky's method which satisfies  $B = TT^T$ . Setting  $v = T^T u$  we have  $P = T^{-1} A (T^{-1})^T$ . This calculation of  $T$  may be accomplished even when  $n$  is large. But in practice the procedure is complicated by the fact that  $P$  need not be sparse even when  $A$  and  $B$  are so. But  $w = Pz$  can be calculated directly from  $A$  and  $T$  without inversion of  $T$ . SIMITZ combined with Cholesky's decomposition eliminates the need for a special inner product - use  $ip(n, w, z) = w^T z$  - and so improves efficiency for large  $n$  or for  $n$  of moderate size. The so called direct methods of treating  $Au = \lambda Bu$  which rely on orthogonal similarity transformations to reduce  $P$  to tridiagonal form followed by application of the QL or QR algorithm [12] are generally unsuitable for large problems unless the structure of  $P$  is quite special. So called iterative methods exemplified by simultaneous iteration must therefore be exploited when  $n$  is large.

One of the most popular of such methods is the inverse power method or inverse iteration [12]. Given an approximate eigenvalue  $\sigma$  and an approximate corresponding eigenvector  $X_0$  the iteration  $(A - \sigma B) X_k = BX_{k-1}$ ,  $k = 1, \dots$  converges rapidly to an eigenvector of  $Au = \lambda Bu$  when at each stage  $X_k^T B X_k = 1$ . This resembles the simultaneous iteration algorithm but when  $p = em = 1$  and when the eigen-



value nearest to  $\sigma$  is sought with its eigenvector. The approximate eigenvalue must be calculated first, however, and inverse iteration may fail if  $\sigma$  tends to identify a multiple eigenvalue. Eigenvalues may be obtained by any of the various bisection methods [4,12] or by Rayleigh-Ritz techniques. The inverse power method and its variants are the only methods applicable to gigantic problems known to this writer.

One of the most promising methods for very large problems is the block Lanczos algorithm [11] for which a FORTRAN program is available to solve the standard symmetric eigenvalue problem  $Pv = \lambda v$ . Block Lanczos has been compared to *ritzit* and found to be superior in many situations. We see no inherent impediment to modifying block Lanczos to produce solutions of  $Au = \lambda Bu$  in the manner in which we have modified *ritzit*.

SECTION IV  
INSTRUCTIONS FOR THE USER

A FORTRAN executable program or subprogram possessing the control statements equivalent to

DIMENSION X(MN,P), D(P), WK(K)

INTEGER P, EM

REAL IP

EXTERNAL IP, INF, OP

may call SUBROUTINE SIMITZ into execution via the statement

CALL SIMITZ (N, P, KM, EPS, IP, OP, INF, EM, X, MN, D, WK)

where

N is an INTEGER input variable, the order of the matrix C.

P is an INTEGER input variable, the number of simultaneous iteration vectors.

KM as an INTEGER input variable is in magnitude the maximum number of iteration steps to be executed. If KM identifies a negative value then P initial approximate eigenvectors are assumed to be present in the array X. Otherwise SIMITZ supplies random initial eigenvectors.

KM as an INTEGER output variable identifies the number KS of iteration steps finally used in the calculation of EM eigenvectors.

EPS is a REAL input variable, the tolerance for accepting eigenvectors.

IP is an EXTERNAL input variable, the name of a FORTRAN compatible REAL FUNCTION subprogram of the form  $IP(N, Z, W)$  which must return the inner product  $(W, BZ) = W^T BZ$  of the vectors identified by the N-arrays Z and W relative to the positive definite matrix B.

OP is an EXTERNAL input variable, the name of a FORTRAN compatible SUBROUTINE subprogram of the form  $OP(N, Z, W)$  which must calculate the image W of the vector identified by the N-array Z under the N-square matrix C without overwriting Z.

INF is an EXTERNAL input variable, the name of a FORTRAN compatible SUBROUTINE subprogram which may be used for obtaining information or to exert control during execution of SIMITZ. INF has the form  $INF(KS, G, H, F)$  where

KS is an INTEGER output variable, the number of the next iteration step.

G is an INTEGER output variable, the number of already accepted eigenvectors.

H is an INTEGER output variable, the number of already accepted eigenvalues.

F is a REAL output variable P-array, error quantities measuring respectively the state of convergence of the P simultaneous iteration vectors. In addition, if convergence fails in SUBROUTINE IMTQL2 after G eigenvectors have been accepted, then  $F(G+1)$  is replaced by  $1000.*FLOAT(IERR)$  where IERR is the error

indicator output by IMTQL2. Each element of the array  
F is initially set by SIMITZ to the value 4.0.

- EM as an INTEGER input variable is the number of eigenvalues to be computed, 0 .LT. EM .LT. P .LE. MN.
- EM as an INTEGER output variable is the number of eigenvectors computed through KM iteration steps.
- X as a real N-by-P input array is a set of P optional initial approximate eigenvectors  $X(I,1), \dots, X(I,P)$ ,  $I = 1, \dots, N$ , interpreted by SIMITZ if KM is negative.
- X as a real N-by-P output array is a set of EM eigenvectors  $X(I,1), \dots, X(I,EM)$ ,  $I = 1, \dots, N$ , computed through IABS(KM) iteration steps with the remainder of X consisting of  $P - EM$  approximate eigenvectors. The N-by-P matrix X satisfies  $X^T B X = I$ , that is, the eigenvectors of C are B-orthonormal.
- MN is an integer input variable which identifies the leading dimension in the calling program of the array X.
- D is a real output P-array of which  $D(1), \dots, D(EM)$  are the eigenvalues of C largest in magnitude in decreasing order corresponding to the computed eigenvectors  $X(I,1), \dots, X(I,EM)$ ,  $I = 1, \dots, N$ .  $D(EM+1), \dots, D(P-1)$  contain approximations to progressively smaller such eigenvalues.  $D(P)$  contains the most recently computed value of E, where the interval  $(-E, E)$  is the interval over which the Chebyshev acceleration was performed.
- WK the initial location of at least  $P^2 + 3P + 2N = K$  consecutive storage locations which may not be overwritten while SIMITZ is



in execution.

SIMITZ employs a DATA statement to assign to a machine dependent REAL variable MT the quotient of the smallest positive REAL value representable by FORTRAN and the smallest REAL value whose sum with 1.0 exceeds 1.0. The performance of SIMITZ is strongly dependent upon the choice of input parameters and upon the careful preparation of the subprograms IP and OP. The user should develop experience with SIMITZ on problems of moderate size before investing processor time on very large problems for which the procedure is ultimately intended. During its execution SIMITZ issues calls to the following subprograms

FUNCTION RANF

returns uniformly distributed random numbers on the open interval (0, 1) given any one argument of any type.

SUBROUTINE TRED2

is the EISPACK program which computes a Householder tridiagonal form of a real symmetric matrix.

SUBROUTINE IMTQL2

is the EISPACK program which computes the eigenvalues and orthonormal eigenvectors of a symmetric tridiagonal matrix.

FUNCTION IP

is described above.

SUBROUTINE OP

is described above.

SUBROUTINE INF

is described above.



AFFDL-TR-76-118

The Function RANF furnished with the CDC FORTRAN Common Library Mathematical Routines is an excellent random number generator. However, no failures of SIMITZ have been noted even when RANF was replaced by a crude in-line congruence procedure. The user should use the USA Standard FORTRAN versions of TRED2 or IMTQL2 rather than the DOUBLE PRECISION versions available for the IBM 360/370 compatible processors. The EISPACK subroutine TQL2 may replace IMTQL2 if desired.

SECTION V

ALGORITHM

A complete listing follows of the simultaneous iteration algorithm for generalized symmetric matrices implemented in USA Standard FORTRAN as SUBROUTINE SIMITZ. The program and its documentation are separately sequenced.

```

SUBROUTINE SIMITZ(N, P, KM, EPS, IP, OP, INF, EM, X, MN, D, WK)  SIMITZ      2
C*****  SIMITZ/0  2
C  IDENTIFICATION  SIMITZ/0  3
C    SIMITZ - ITERATIVE COMPUTATION OF EIGENVALUES LARGEST IN MAGNI-  SIMITZ/0  4
C    TUDE AND CORRESPONDING EIGENVECTORS OF A REAL GENERAL-  SIMITZ/0  5
C    IZED SYMMETRIC MATRIX  SIMITZ/0  6
C  FORTRAN SUBROUTINE SUBPROGRAM  SIMITZ/0  7
C  US AIR FORCE FLIGHT DYNAMICS LABORATORY  SIMITZ/0  8
C  WRIGHT-PATTERSON AFB, OHIO 45433  SIMITZ/0  9
C  PURPOSE  SIMITZ/0 10
C  A REAL N-SQUARE MATRIX C IS B-SYMMETRIC RELATIVE TO AN N-SQUARE  SIMITZ/0 11
C  POSITIVE DEFINITE MATRIX B IN CASE  $BC = C'B$  WHERE  $C'$  IS THE  SIMITZ/0 12
C  TRANSPOSE OF C. GIVEN AN OPTIONAL SET OF P INITIAL APPROXIMATE  SIMITZ/0 13
C  EIGENVECTORS OF A REAL N-SQUARE B-SYMMETRIC MATRIX C CORRES-  SIMITZ/0 14
C  PONDING TO P EIGENVALUES OF C LARGEST IN MAGNITUDE, SIMITZ COM-  SIMITZ/0 15
C  PUTES EM EIGENVALUES AND EM CORRESPONDING EIGENVECTORS TO A  SIMITZ/0 16
C  PRECISION DEPENDENT ON THE STRUCTURE OF B AND C AND ON A GIVEN  SIMITZ/0 17
C  TOLERANCE EPS. THE MATRIX B IS PRESENTED TO SIMITZ AS AN ALGO-  SIMITZ/0 18
C  RITHM FOR CALCULATING THE STANDARD INNER PRODUCT  $(W, BZ) = W'BZ$   SIMITZ/0 19
C  GIVEN COLUMN N-VECTORS W AND Z IMPLEMENTED AS A FORTRAN COM-  SIMITZ/0 20
C  PATIBLE REAL FUNCTION SUBPROGRAM. THE MATRIX C IS PRESENTED AS  SIMITZ/0 21
C  A SUBROUTINE SUBPROGRAM WHICH GIVEN A COLUMN N-VECTOR Z CALCU-  SIMITZ/0 22
C  LATES ITS IMAGE  $W = CZ$  UNDER THE B-SYMMETRIC MATRIX C. DEPEND-  SIMITZ/0 23
C  ING ON THE CHOICE OF B AND C, SIMITZ APPLIES TO A WIDE VARIETY  SIMITZ/0 24
C  OF SYMMETRIC EIGENPROBLEMS.  SIMITZ/0 25
C  CONTROL  SIMITZ/0 26
C  DIMENSION X(MN,P), D(P), WK(K)  SIMITZ/0 27
C  INTEGER P, EM  SIMITZ/0 28
C  REAL IP  SIMITZ/0 29
C  EXTERNAL IP, INF, OP  SIMITZ/0 30
C  .  SIMITZ/0 31
C  .  SIMITZ/0 32
C  .  SIMITZ/0 33
C  .  SIMITZ/0 34
C  CALL SIMITZ(N, P, KM, EPS, IP, OP, INF, EM, X, MN, D, WK)  SIMITZ/0 35
C  SIMITZ/0 36
C  WHERE  SIMITZ/0 37
C  N IS AN INTEGER INPUT VARIABLE, THE ORDER OF THE MATRIX C.  SIMITZ/0 38
C  P IS AN INTEGER INPUT VARIABLE, THE NUMBER OF SIMULTANEOUS  SIMITZ/0 39
C  ITERATION VECTORS.  SIMITZ/0 40
C  KM AS AN INTEGER INPUT VARIABLE IS IN MAGNITUDE THE MAXIMUM  SIMITZ/0 41
C  NUMBER OF ITERATION STEPS TO BE EXECUTED. IF KM IDENTIFIES  SIMITZ/0 42
C  A NEGATIVE VALUE THEN P INITIAL APPROXIMATE EIGENVECTORS  SIMITZ/0 43
C  ARE ASSUMED TO BE PRESENT IN THE ARRAY X. OTHERWISE SIMITZ  SIMITZ/0 44
C  SUPPLIES RANDOM INITIAL EIGENVECTORS.  SIMITZ/0 45
C  KM AS AN INTEGER OUTPUT VARIABLE IDENTIFIES THE NUMBER KS OF  SIMITZ/0 46
C  ITERATION STEPS FINALLY USED IN THE CALCULATION OF EM  SIMITZ/0 47
C  EIGENVECTORS.  SIMITZ/0 48
C  EPS IS A REAL INPUT VARIABLE, THE TOLERANCE FOR ACCEPTING  SIMITZ/0 49
C  EIGENVECTORS.  SIMITZ/0 50
C  IP IS AN EXTERNAL INPUT VARIABLE, THE NAME OF A FORTRAN COM-  SIMITZ/0 51
C  PATIBLE REAL FUNCTION SUBPROGRAM OF THE FORM IP(V, Z, W)  SIMITZ/0 52
C  WHICH MUST RETURN THE INNER PRODUCT  $(W, BZ) = W'BZ$  OF THE  SIMITZ/0 53
C  VECTORS IDENTIFIED BY THE N-ARRAYS Z AND W RELATIVE TO THE  SIMITZ/0 54
C  POSITIVE DEFINITE MATRIX B.  SIMITZ/0 55
C  OP IS AN EXTERNAL INPUT VARIABLE, THE NAME OF A FORTRAN COM-  SIMITZ/0 56
C  PATIBLE SUBROUTINE SUBPROGRAM OF THE FORM OP(N, Z, W)  SIMITZ/0 57

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WHICH MUST CALCULATE THE IMAGE W OF THE VECTOR IDENTIFIED BY THE N-ARRAY Z UNDER THE N-SQUARE MATRIX C WITHOUT OVERWRITING Z.  
 INF IS AN EXTERNAL INPUT VARIABLE, THE NAME OF A FORTRAN COMPATIBLE SUBROUTINE SUBPROGRAM WHICH MAY BE USED FOR OBTAINING INFORMATION OR TO EXERT CONTROL DURING EXECUTION OF SIMITZ. INF HAS THE FORM INF(KS, G, H, F) WHERE  
 KS IS AN INTEGER OUTPUT VARIABLE, THE NUMBER OF THE NEXT ITERATION STEP.  
 G IS AN INTEGER OUTPUT VARIABLE, THE NUMBER OF ALREADY ACCEPTED EIGENVECTORS.  
 H IS AN INTEGER OUTPUT VARIABLE, THE NUMBER OF ALREADY ACCEPTED EIGENVALUES.  
 F IS A REAL OUTPUT VARIABLE P-ARRAY, ERROR QUANTITIES MEASURING RESPECTIVELY THE STATE OF CONVERGENCE OF THE P SIMULTANEOUS ITERATION VECTORS. IN ADDITION, IF CONVERGENCE FAILS IN SUBROUTINE INTOL2 AFTER G EIGENVECTORS HAVE BEEN ACCEPTED, THEN F(G+1) IS REPLACED BY 1000.\*FLOAT(IERR) WHERE IERR IS THE ERROR INDICATOR OUTPUT BY INTOL2. EACH ELEMENT OF THE ARRAY F IS INITIALLY SET BY SIMITZ TO THE VALUE 4.0.  
 EM AS AN INTEGER INPUT VARIABLE IS THE NUMBER OF EIGENVALUES TO BE COMPUTED, 0 .LT. EM .LT. P .LE. N .LE. MN.  
 EM AS AN INTEGER OUTPUT VARIABLE IS THE NUMBER OF EIGENVECTORS COMPUTED THROUGH KM ITERATION STEPS.  
 X AS A REAL N-BY-P INPUT ARRAY IS A SET OF P OPTIONAL INITIAL APPROXIMATE EIGENVECTORS X(I,1), ..., X(I,P), I = 1, ..., N, INTERPRETED BY SIMITZ IF KM IS NEGATIVE.  
 X AS A REAL N-BY-P OUTPUT ARRAY IS A SET OF EM EIGENVECTORS X(I,1), ..., X(I,EM), I = 1, ..., N, COMPUTED THROUGH IABS(KM) ITERATION STEPS WITH THE REMAINDER OF X CONSISTING OF P - EM APPROXIMATE EIGENVECTORS. THE N-BY-P MATRIX X SATISFIES  $X^*BX = I$ , THAT IS, THE EIGENVECTORS OF C ARE B-ORTHOGONAL.  
 MN IS AN INTEGER INPUT VARIABLE WHICH IDENTIFIES THE LEADING DIMENSION IN THE CALLING PROGRAM OF THE ARRAY X.  
 D IS A REAL OUTPUT P-ARRAY OF WHICH D(1), ..., D(EM) ARE THE EIGENVALUES OF C LARGEST IN MAGNITUDE IN DECREASING ORDER CORRESPONDING TO THE COMPUTED EIGENVECTORS X(I,1), ..., X(I,EM), I = 1, ..., N. D(EM+1), ..., D(P-1) CONTAIN APPROXIMATIONS TO PROGRESSIVELY SMALLER SUCH EIGENVALUES. D(P) CONTAINS THE MOST RECENTLY COMPUTED VALUE OF E, WHERE THE INTERVAL (-E, E) IS THE INTERVAL OVER WHICH THE CHEBYSHEV ACCELERATION WAS PERFORMED.  
 WK THE INITIAL LOCATION OF AT LEAST  $P^2 + 3P + 2N = K$  CONSECUTIVE STORAGE LOCATIONS WHICH MAY NOT BE OVERWRITTEN WHILE SIMITZ IS IN EXECUTION.  
 OTHER PROGRAMMING INFORMATION  
 SIMITZ EMPLOYS A DATA STATEMENT TO ASSIGN TO A MACHINE DEPENDENT REAL VARIABLE MT THE QUOTIENT OF THE SMALLEST POSITIVE REAL VALUE REPRESENTABLE BY FORTRAN AND THE SMALLEST REAL VALUE WHOSE SUM WITH 1.0 EXCEEDS 1.0.  
 THE PERFORMANCE OF SIMITZ IS STRONGLY DEPENDENT UPON THE CHOICE OF INPUT PARAMETERS AND UPON THE CAREFUL PREPARATION OF THE SUBPROGRAMS IP AND OP. THE USER SHOULD DEVELOP EXPERIENCE WITH SIMITZ ON PROBLEMS OF MODERATE SIZE BEFORE INVESTING



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C      PROCESSOR TIME ON VERY LARGE PROBLEMS FOR WHICH THE PROCEDURE SIMITZ/D115
C      IS ULTIMATELY INTENDED. SIMITZ/D116
C      OTHER PROGRAMS REQUIRED SIMITZ/D117
C      FUNCTION RANF SIMITZ/D118
C      RETURNS UNIFORMLY DISTRIBUTED RANDOM NUMBERS ON THE OPEN SIMITZ/D119
C      INTERVAL (0, 1) GIVEN ANY ONE ARGUMENT OF ANY TYPE. SIMITZ/D120
C      SUBROUTINE TRED2 SIMITZ/D121
C      IS THE EISPACK (4) PROGRAM WHICH COMPUTES A HOUSEHOLDER SIMITZ/D122
C      TRIDIAGONAL FORM OF A REAL SYMMETRIC MATRIX. SIMITZ/D123
C      SUBROUTINE IMTQL2 SIMITZ/D124
C      IS THE EISPACK PROGRAM WHICH COMPUTES THE EIGENVALUES AND SIMITZ/D125
C      ORTHONORMAL EIGENVECTORS OF A SYMMETRIC TRIDIAGONAL MATRIX. SIMITZ/D126
C      FUNCTION IP SIMITZ/D127
C      IS DESCRIBED ABOVE. SIMITZ/D128
C      SUBROUTINE OP SIMITZ/D129
C      IS DESCRIBED ABOVE. SIMITZ/D130
C      SUBROUTINE INF SIMITZ/D131
C      IS DESCRIBED ABOVE. SIMITZ/D132
C      METHOD SIMITZ/D133
C      SIMITZ REPRESENTS RESULTS OF EXTENSIVE MODIFICATIONS AND TESTS SIMITZ/D134
C      OF SUBROUTINE RITZIT (1), AN ANSI FORTRAN TRANSLATION OF THE SIMITZ/D135
C      ALGOL 60 PROCEDURE OF THE SAME NAME (3). THE BASIC RUTISHAUSER SIMITZ/D136
C      -REINSCH ALGORITHM IS PRESERVED. SIMITZ/D137
C      REFERENCES SIMITZ/D138
C      (1) PAUL J. NIKOLAI AND NAI-KUAN TSAO, THE AFL LINEAR ALGEBRA SIMITZ/D139
C      LIBRARY HANDBOOK, AFL TR 74-0106, AEROSPACE RESEARCH LABOR- SIMITZ/D140
C      ATORIES, WRIGHT-PATTERSON AFB, OHIO, 1974. SIMITZ/D141
C      (2) HEINZ RUTISHAUSER, COMPUTATIONAL ASPECTS OF F.L. BAUER'S SIMITZ/D142
C      SIMULTANEOUS ITERATION METHOD, NUMER. MATH. 13(1969), 4-13. SIMITZ/D143
C      (3) -----, SIMULTANEOUS ITERATION METHOD FOR SYM- SIMITZ/D144
C      METRIC MATRICES, NUMER. MATH. 16(1970), 205-223. SIMITZ/D145
C      (4) R.T. SMITH ET AL, MATRIX EIGENSYSTEM ROUTINES-EISPACK SIMITZ/D146
C      GUIDE, LECTURE NOTES IN COMPUTER SCIENCE 5, SPRINGER-VERLAG SIMITZ/D147
C      NEW YORK, 1974. SIMITZ/D148
C      ***** SIMITZ/D149
C      EXTERNAL INF, IP, OP SIMITZ 4
C      INTEGER EM, G, H, I, IG, IK, J, SIMITZ 5
C      * JK, JP, K, KM, KS, L, LF, SIMITZ 6
C      * L1, M, MN, M1, N, P, Z1, SIMITZ 7
C      * Z2 SIMITZ 8
C      LOGICAL ORIG SIMITZ 9
C      REAL D, E, EPS, E1, E2, IP, MT, SIMITZ 10
C      * S, T, WK, X SIMITZ 11
C      DIMENSION X(MN,1), D(1), WK(P,P,1) SIMITZ 12
C      DATA MT / .220360641585062E-279/ SIMITZ 13
C      THE LOCAL VARIABLE ARRAYS FROM (3) ARE ASSIGNED TO THE SIMITZ 14
C      VARIABLE ARRAY WK AS FOLLOWS: SIMITZ 15
C      WK(I,J,1) = R(I,J), I, J = 1, ..., P. SIMITZ 16
C      WK(I,1,2) = CX(I), I = 1, ..., P. SIMITZ 17
C      WK(I,2,2) = F(I), I = 1, ..., P. SIMITZ 18
C      WK(I,3,2) = RO(I), I = 1, ..., P. SIMITZ 19
C      WK(I,4,2) = U(I), I = 1, ..., N. SIMITZ 20
C      WK(I+N,4,2) = W(I), I = 1, ..., N. SIMITZ 21
C      NOT NEEDED ARE V(I), I = 1, ..., N, R(I), I = 1, ..., P, AND SIMITZ 22
C      SIMITZ 23
C      SIMITZ 24
C      SIMITZ 25

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C	Q(I,J), I, J = 1, ..., P.	SIMITZ	26
C		SIMITZ	27
C	THE NEXT STATEMENT IS START.	SIMITZ	28
	E = .0E+00	SIMITZ	29
	G = 0	SIMITZ	30
	IG = 1	SIMITZ	31
	H = 0	SIMITZ	32
	Z1 = 0	SIMITZ	33
	Z2 = 0	SIMITZ	34
	KS = 0	SIMITZ	35
	M = 1	SIMITZ	36
	WK(P,1,2) = .0E+00	SIMITZ	37
	DO 10 L = 1, P	SIMITZ	38
	WK(L,2,2) = .4E+01	SIMITZ	39
	WK(L,3,2) = .0E+00	SIMITZ	40
10	CONTINUE	SIMITZ	41
	IF (KM) 50, 50, 20	SIMITZ	42
20	DO 40 L = 1, P	SIMITZ	43
	DO 30 J = 1, N	SIMITZ	44
	X(J,L) = .2E+01*RANF(.2E+01) - .1E+01	SIMITZ	45
30	CONTINUE	SIMITZ	46
40	CONTINUE	SIMITZ	47
50	KM = IABS(KM)	SIMITZ	48
	ASSIGN 60 TO IK	SIMITZ	49
	LF = IG	SIMITZ	50
	L1 = P	SIMITZ	51
	GO TO 990	SIMITZ	52
C	RAYLEIGH-RITZ STEP	SIMITZ	53
C	STATEMENT 60 IS LOOP.	SIMITZ	54
60	DO 80 K = IG, P	SIMITZ	55
	CALL OP(N, X(1,K), WK(1,4,2))	SIMITZ	56
	DO 70 J = 1, N	SIMITZ	57
	X(J,K) = WK(J,4,2)	SIMITZ	58
70	CONTINUE	SIMITZ	59
80	CONTINUE	SIMITZ	60
	ASSIGN 90 TO IK	SIMITZ	61
	LF = IG	SIMITZ	62
	L1 = P	SIMITZ	63
	GO TO 990	SIMITZ	64
90	IF (KS) 150, 100, 150	SIMITZ	65
C	MEASURES AGAINST UNHAPPY CHOICE OF INITIAL VECTORS	SIMITZ	66
100	DO 130 K = 1, P	SIMITZ	67
	IF (WK(K,K,1)) 130, 110, 130	SIMITZ	68
110	DO 120 I = 1, N	SIMITZ	69
	X(I,K) = .2E+01*RANF(.2E+01) - .1E+01	SIMITZ	70
120	CONTINUE	SIMITZ	71
	KS = 1	SIMITZ	72
130	CONTINUE	SIMITZ	73
	IF (KS - 1) 150, 140, 150	SIMITZ	74
140	ASSIGN 60 TO IK	SIMITZ	75
	LF = 1	SIMITZ	76
	L1 = P	SIMITZ	77
	GO TO 990	SIMITZ	78
150	DO 180 K = IG, P	SIMITZ	79
	DO 170 L = K, P	SIMITZ	80
	S = .0E+00	SIMITZ	81
	DO 160 I = L, P	SIMITZ	82

	S = S + WK(I,K,1)*WK(I,L,1)	SIMITZ	83
160	CONTINUE	SIMITZ	84
	WK(L,K,1) = -S	SIMITZ	85
170	CONTINUE	SIMITZ	86
190	CONTINUE	SIMITZ	87
	CALL TRED2(P, P = G, WK(IG,IG,1), D(IG), WK(1,4,2), WK(IG,IG,1))	SIMITZ	88
	CALL INTOL2(P, P = G, D(IG), WK(1,4,2), WK(IG,IG,1), L)	SIMITZ	89
	WK(IG,2,2) = AMAX1(WK(IG,2,2), .1E+04*FLOAT(L))	SIMITZ	90
	DO 193 K = IG, P	SIMITZ	91
	D(K) = SQRT(AMAX1(-D(K), .0E+00))	SIMITZ	92
190	CONTINUE	SIMITZ	93
C	REORDERING EIGENVALUES AND EIGENVECTORS ACCORDING TO SIZE OF	SIMITZ	94
C	THE FORMER IS ACCOMPLISHED IN SUBROUTINE INTOL2.	SIMITZ	95
	DO 230 J = 1, N	SIMITZ	96
	DO 210 K = IG, P	SIMITZ	97
	S = .JE+00	SIMITZ	98
	DO 200 L = IG, P	SIMITZ	99
	S = S + X(J,L)*WK(L,K,1)	SIMITZ	100
200	CONTINUE	SIMITZ	101
	WK(K,4,2) = S	SIMITZ	102
210	CONTINUE	SIMITZ	103
	DO 220 K = IG, P	SIMITZ	104
	X(J,K) = WK(K,4,2)	SIMITZ	105
220	CONTINUE	SIMITZ	106
230	CONTINUE	SIMITZ	107
	KS = KS + 1	SIMITZ	108
	E = AMAX1(D(P), E)	SIMITZ	109
C	RANDOMIZATION	SIMITZ	110
	IF (3 - 71) 260, 240, 240	SIMITZ	111
240	DO 250 J = 1, N	SIMITZ	112
	X(J,P) = .2E+01*RAVF(.2E+01) - .1E+01	SIMITZ	113
250	CONTINUE	SIMITZ	114
	JP = P - 1	SIMITZ	115
	ASSIGN 260 TO IK	SIMITZ	116
	LF = P	SIMITZ	117
	L1 = P	SIMITZ	118
	GO TO 990	SIMITZ	119
C	COMPUTE CONTROL QUANTITIES CX(I).	SIMITZ	120
260	DO 310 K = IG, JP	SIMITZ	121
	S = (D(K) - E)*(D(K) + E)	SIMITZ	122
	IF (S) 270, 270, 280	SIMITZ	123
270	WK(K,1,2) = .0E+00	SIMITZ	124
	GO TO 310	SIMITZ	125
280	IF (E) 300, 290, 300	SIMITZ	126
290	WK(K,1,2) = .1E+04 + ALOG(D(K))	SIMITZ	127
	GO TO 310	SIMITZ	128
300	WK(K,1,2) = ALOG((D(K) + SQRT(S))/E)	SIMITZ	129
310	CONTINUE	SIMITZ	130
C	ACCEPTANCE TEST FOR EIGENVALUES INCLUDING ADJUSTMENT OF EM AND	SIMITZ	131
C	H SUCH THAT D(EM) .GT. E, D(H) .GT. E AND D(EM) DOES NOT	SIMITZ	132
C	OSCILLATE STRONGLY	SIMITZ	133
	I = 71 - 1	SIMITZ	134
	K = G	SIMITZ	135
320	K = K + 1	SIMITZ	136
	IF (EM - K) 370, 330, 330	SIMITZ	137
330	IF (D(K) - E) 360, 360, 340	SIMITZ	138
340	IF (I) 320, 320, 350	SIMITZ	139

350	IF (D(K) - .999E+00*WK(K,3,2)) 360, 360, 320	SIMITZ	140
360	CONTINUE	SIMITZ	141
	EM = K - 1	SIMITZ	142
C	STATEMENT 370 IS EX4.	SIMITZ	143
370	IF (EM) 380, 1130, 380	SIMITZ	144
380	K = H	SIMITZ	145
	S = .1E+01 + .1E+00*EPS	SIMITZ	146
390	K = K + 1	SIMITZ	147
	IF (D(K)) 400, 410, 400	SIMITZ	148
400	IF (D(K) - S*WK(K,3,2)) 390, 390, 410	SIMITZ	149
410	CONTINUE	SIMITZ	150
	H = K - 1	SIMITZ	151
	K = EM	SIMITZ	152
420	K = K + 1	SIMITZ	153
	IF (K - H) 430, 430, 450	SIMITZ	154
430	IF (D(K) - E) 440, 440, 420	SIMITZ	155
440	CONTINUE	SIMITZ	156
	H = K - 1	SIMITZ	157
C	ACCEPTANCE TEST FOR EIGENVECTORS	SIMITZ	158
450	L = G	SIMITZ	159
	E2 = .0E+00	SIMITZ	160
	DO 590 K = IG, JP	SIMITZ	161
	IF (K - (L + 1)) 510, 460, 510	SIMITZ	162
C	CHECK FOR NESTED EIGENVALUES	SIMITZ	163
460	L = K	SIMITZ	164
	L1 = K	SIMITZ	165
	S = .5E+00/FLOAT(KS)	SIMITZ	166
	T = .1E+01/FLOAT(KS*M)	SIMITZ	167
470	L = L + 1	SIMITZ	168
	IF (L - JP) 480, 480, 490	SIMITZ	169
480	IF (WK(L,1,2)*(WK(L,1,2) + S) + T - WK(L-1,1,2)*WK(L-1,1,2))	SIMITZ	170
	490, 490, 470	SIMITZ	171
490	CONTINUE	SIMITZ	172
	L = L - 1	SIMITZ	173
C	THE NEXT STATEMENT IS EX5.	SIMITZ	174
	IF (L - H) 510, 510, 500	SIMITZ	175
500	L = L1 - 1	SIMITZ	176
	GO TO 600	SIMITZ	177
510	CALL OP(N, X(1,K), WK(1,4,2))	SIMITZ	178
	S = .0E+00	SIMITZ	179
	DO 540 J = 1, L	SIMITZ	180
	IF (ABS(D(J) - D(K)) - .1E-01*D(K)) 520, 540, 540	SIMITZ	181
520	T = IP(N, WK(1,4,2), X(1,J))	SIMITZ	182
	DO 530 I = 1, N	SIMITZ	183
	WK(I,4,2) = WK(I,4,2) - T*X(I,J)	SIMITZ	184
530	CONTINUE	SIMITZ	185
	S = S + T*T	SIMITZ	186
540	CONTINUE	SIMITZ	187
	T = .1E+01	SIMITZ	188
	IF (S .NE. .0E+00) T = IP(N, WK(1,4,2), WK(1,4,2))	SIMITZ	189
	E2 = AMAX1(E2, SORT(T/(S + T)))	SIMITZ	190
	IF (K - L) 590, 550, 590	SIMITZ	191
C	TEST FOR ACCEPTANCE OF GROUP OF EIGENVECTORS	SIMITZ	192
550	IF (L .GE. EM .AND. D(EM)*WK(EM,2,2) .LT. EPS*(D(EM) - E))	SIMITZ	193
	G = EM	SIMITZ	194
	IF (E2 - WK(L,2,2)) 560, 580, 580	SIMITZ	195
560	DO 570 J = L1, L	SIMITZ	196



WK(J,2,2) = E2	SIMITZ 197
570 CONTINUE	SIMITZ 198
580 IF (L .LE. EM .AND. D(L)*WK(L,2,2) .LT. EPS*(D(L) - E)) G = L	SIMITZ 199
590 CONTINUE	SIMITZ 200
C ADJUST M.	SIMITZ 201
C STATEMENT 600 IS EX6.	SIMITZ 202
600 IG = G + 1	SIMITZ 203
IF (E - .4E-01*D(1)) 610, 610, 620	SIMITZ 204
610 M = 1	SIMITZ 205
K = 1	SIMITZ 206
GO TO 630	SIMITZ 207
620 E2 = .2E+01/E	SIMITZ 208
E1 = .51E+00*E2	SIMITZ 209
K = 2*INT(.4E+01/AMIN1(WK(1,1,2), .4E+(1)))	SIMITZ 210
M = MIN0(M, K)	SIMITZ 211
C REDUCE EM IF CONVERGENCE WOULD BE TOO SLOW.	SIMITZ 212
630 IF (WK(EM,2,2)) 640, 690, 640	SIMITZ 213
640 IF (FLOAT(KS) - .9E+00*FLOAT(KM)) 650, 690, 690	SIMITZ 214
650 S = FLOAT(K)*WK(EM,1,2)	SIMITZ 215
IF (S - .5E-01) 660, 670, 670	SIMITZ 216
660 T = .5E+00*S*WK(EM,1,2)	SIMITZ 217
GO TO 680	SIMITZ 218
670 T = WK(EM,1,2) + ALOG(.5E+00 + .5E+00*EXP(-.2E+01*S))/FLOAT(K)	SIMITZ 219
680 S = ALOG(D(EM)*WK(EM,2,2)/(EPS*(D(EM) - E)))	SIMITZ 220
IF (S*FLOAT(KS) .GT. T*FLOAT((KM - KS)*KM)) EM = EM - 1	SIMITZ 221
C STATEMENT 690 IS EX2.	SIMITZ 222
690 DO 700 K = IG, JP	SIMITZ 223
WK(K,3,2) = D(K)	SIMITZ 224
700 CONTINUE	SIMITZ 225
CALL INF(KS, G, H, WK(1,2,2))	SIMITZ 226
IF (G .GE. EM .OR. KS .GE. KM) GO TO 1130	SIMITZ 227
C STATEMENT 710 IS EX1.	SIMITZ 228
710 IF (KS + M - KM) 730, 730, 720	SIMITZ 229
720 Z2 = -1	SIMITZ 230
IF (M .GT. 1) M = 2*((KM - KS + 1)/2)	SIMITZ 231
730 M1 = M	SIMITZ 232
C SHORTCUT LAST INTERMEDIATE BLOCK IF ALL F(I) ARE SUFFICIENTLY	SIMITZ 233
C SMALL.	SIMITZ 234
IF (L - EM) 740, 740, 740	SIMITZ 235
740 S = D(EM)*WK(EM,2,2)/(EPS*(D(EM) - E))	SIMITZ 236
T = S*S - .1E+01	SIMITZ 237
IF (T) 60, 60, 750	SIMITZ 238
750 S = ALOG(S + SQRT(T))/(WK(EM,1,2) - WK(H+1,1,2))	SIMITZ 239
M1 = 2*INT(.5E+00*S + .101E+01)	SIMITZ 240
IF (M1 - M) 770, 770, 760	SIMITZ 241
760 M1 = M	SIMITZ 242
GO TO 780	SIMITZ 243
770 Z2 = -1	SIMITZ 244
C CHERYSHEV ITERATION	SIMITZ 245
780 IF (M - 1) 800, 790, 820	SIMITZ 246
790 DO 810 K = IG, P	SIMITZ 247
CALL OP(N, X(1,K), WK(1,4,2))	SIMITZ 248
DO 800 I = 1, N	SIMITZ 249
X(I,K) = WK(I,4,2)	SIMITZ 250
800 CONTINUE	SIMITZ 251
810 CONTINUE	SIMITZ 252
GO TO 900	SIMITZ 253

820	L1 = M1 - 4	SIMITZ	254
	DO 890 K = IG, P	SIMITZ	255
	CALL OP(N, X(1,K), WK(1,4,2))	SIMITZ	256
	DO 830 I = 1, N	SIMITZ	257
	IK = I + N	SIMITZ	258
	WK(IK,4,2) = E1*WK(I,4,2)	SIMITZ	259
830	CONTINUE	SIMITZ	260
	CALL OP(N, WK(N+1,4,2), WK(1,4,2))	SIMITZ	261
	DO 840 I = 1, N	SIMITZ	262
	X(I,K) = E2*WK(I,4,2) - X(I,K)	SIMITZ	263
840	CONTINUE	SIMITZ	264
	IF (L1) 890, 850, 850	SIMITZ	265
850	DO 860 J = 4, M1, 2	SIMITZ	266
	CALL OP(N, X(1,K), WK(1,4,2))	SIMITZ	267
	DO 860 I = 1, N	SIMITZ	268
	IK = I + N	SIMITZ	269
	WK(IK,4,2) = E2*WK(I,4,2) - WK(IK,4,2)	SIMITZ	270
860	CONTINUE	SIMITZ	271
	CALL OP(N, WK(N+1,4,2), WK(1,4,2))	SIMITZ	272
	DO 870 I = 1, N	SIMITZ	273
	X(I,K) = E2*WK(I,4,2) - X(I,K)	SIMITZ	274
870	CONTINUE	SIMITZ	275
880	CONTINUE	SIMITZ	276
890	CONTINUE	SIMITZ	277
900	ASSIGN 910 TO IK	SIMITZ	278
	LF = IG	SIMITZ	279
	L1 = P	SIMITZ	280
	GO TO 990	SIMITZ	281
C	DISCOUNTING THE ERROR QUANTITIES F	SIMITZ	282
910	IF (G - H) 920, 970, 970	SIMITZ	283
920	IF (M - 1) 950, 930, 950	SIMITZ	284
930	DO 940 K = IG, H	SIMITZ	285
	WK(K,2,2) = WK(K,2,2)*(D(H+1)/D(K))	SIMITZ	286
940	CONTINUE	SIMITZ	287
	GO TO 970	SIMITZ	288
950	T = EXP(-FLOAT(M1)*WK(H+1,1,2))	SIMITZ	289
	DO 960 K = IG, H	SIMITZ	290
	S = EXP(-FLOAT(M1)*(WK(K,1,2) - WK(H+1,1,2)))	SIMITZ	291
	WK(K,2,2) = S*WK(K,2,2)*(0.1E+01 + T*T)/(0.1E+01 + (S*T)*(S*T))	SIMITZ	292
960	CONTINUE	SIMITZ	293
970	KS = KS + M1	SIMITZ	294
	Z2 = Z2 - M1	SIMITZ	295
C	POSSIBLE REPETITION OF INTERMEDIATE STEPS	SIMITZ	296
	IF (Z2) 980, 710, 710	SIMITZ	297
980	Z1 = Z1 + 1	SIMITZ	298
	Z2 = 2*Z1	SIMITZ	299
	M = M + M	SIMITZ	300
	GO TO 60	SIMITZ	301
C	PERFORMS ORTHONORMALIZATION OF COLUMNS 1 THROUGH L1 OF ARRAY	SIMITZ	302
C	X ASSUMING THAT COLUMNS 1 THROUGH LF - 1 ARE ALREADY ORTHO-	SIMITZ	303
C	NORMAL	SIMITZ	304
990	DO 1120 K = LF, L1	SIMITZ	305
	ORIG = .TRUE.	SIMITZ	306
1000	T = .0E+00	SIMITZ	307
	JK = K - 1	SIMITZ	308
	IF (JK) 1040, 1040, 1010	SIMITZ	309
1010	DO 1030 I = 1, JK	SIMITZ	310

S = IP(N, X(1,I), X(1,K))	SIMITZ 311
IF (ORIG) WK(K,I,1) = S	SIMITZ 312
T = T + S*S	SIMITZ 313
DO 1020 J = 1, N	SIMITZ 314
X(J,K) = X(J,K) - S*X(J,I)	SIMITZ 315
1020 CONTINUE	SIMITZ 316
1030 CONTINUE	SIMITZ 317
1040 S = IP(N, X(1,K), X(1,K))	SIMITZ 318
T = S + T	SIMITZ 319
IF (S - .1E-01*T) 1060, 1060, 1050	SIMITZ 320
1050 IF (T - MT) 1060, 1060, 1060	SIMITZ 321
1060 ORIG = .FALSE.	SIMITZ 322
IF (S - MT) 1070, 1070, 1000	SIMITZ 323
1070 S = .0E+00	SIMITZ 324
1080 S = SORT(S)	SIMITZ 325
WK(K,K,1) = S	SIMITZ 326
IF (S) 1090, 1100, 1090	SIMITZ 327
1090 S = .1E+01/S	SIMITZ 328
1100 DO 1110 J = 1, N	SIMITZ 329
X(J,K) = S*X(J,K)	SIMITZ 330
1110 CONTINUE	SIMITZ 331
1120 CONTINUE	SIMITZ 332
GO TO IK, (60, 90, 250, 910, 1140)	SIMITZ 333
C STATEMENT 1130 IS EX.	SIMITZ 334
1130 EM = G	SIMITZ 335
C SOLVE EIGENVALUE PROBLEM OF PROJECTION OF MATRIX G.	SIMITZ 336
ASSIGN 1140 TO IK	SIMITZ 337
LF = 1	SIMITZ 338
L1 = JP	SIMITZ 339
GO TO 990	SIMITZ 340
1140 DO 1160 K = 1, JP	SIMITZ 341
CALL OP(N, X(1,K), X(1,P))	SIMITZ 342
DO 1150 I = 1, K	SIMITZ 343
WK(K,I,1) = -IP(N, X(1,I), X(1,P))	SIMITZ 344
1150 CONTINUE	SIMITZ 345
1160 CONTINUE	SIMITZ 346
CALL TRED2(P, JP, WK, 0, WK(1,4,2), WK)	SIMITZ 347
CALL IMTOL2(P, JP, 0, WK(1,4,2), WK, L)	SIMITZ 348
WK(IG,2,2) = AMAX1(WK(IG,2,2), .1E+04*FLOAT(L))	SIMITZ 349
C ARRANGE EIGENVALUES IN ORDER OF DECREASING ABSOLUTE VALUE.	SIMITZ 350
DO 1210 J = 1, JP	SIMITZ 351
K = J	SIMITZ 352
DO 1170 I = J, JP	SIMITZ 353
IF (ABS(D(I)) .GT. ABS(D(K))) K = I	SIMITZ 354
1170 CONTINUE	SIMITZ 355
IF (K - J) 1200, 1200, 1180	SIMITZ 356
1180 T = D(K)	SIMITZ 357
D(K) = D(J)	SIMITZ 358
D(J) = T	SIMITZ 359
DO 1190 I = 1, JP	SIMITZ 360
T = WK(I,K,1)	SIMITZ 361
WK(I,K,1) = WK(I,J,1)	SIMITZ 362
WK(I,J,1) = T	SIMITZ 363
1190 CONTINUE	SIMITZ 364
1200 D(J) = -D(J)	SIMITZ 365
1210 CONTINUE	SIMITZ 366
DO 1250 J = 1, N	SIMITZ 367

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DO 1230 I = 1, JP
  S = .0E+00
  DO 1220 K = 1, JP
    S = S + X(J,K)*WK(K,I,1)
1220  CONTINUE
    WK(I,4,2) = S
1230  CONTINUE
    DO 1240 I = 1, JP
      X(J,I) = WK(I,4,2)
1240  CONTINUE
1250  CONTINUE
    KM = KS
    O(P) = E
    RETURN
    END

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SIMITZ 368
SIMITZ 369
SIMITZ 370
SIMITZ 371
SIMITZ 372
SIMITZ 373
SIMITZ 374
SIMITZ 375
SIMITZ 376
SIMITZ 377
SIMITZ 378
SIMITZ 379
SIMITZ 380
SIMITZ 381
SIMITZ 382

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